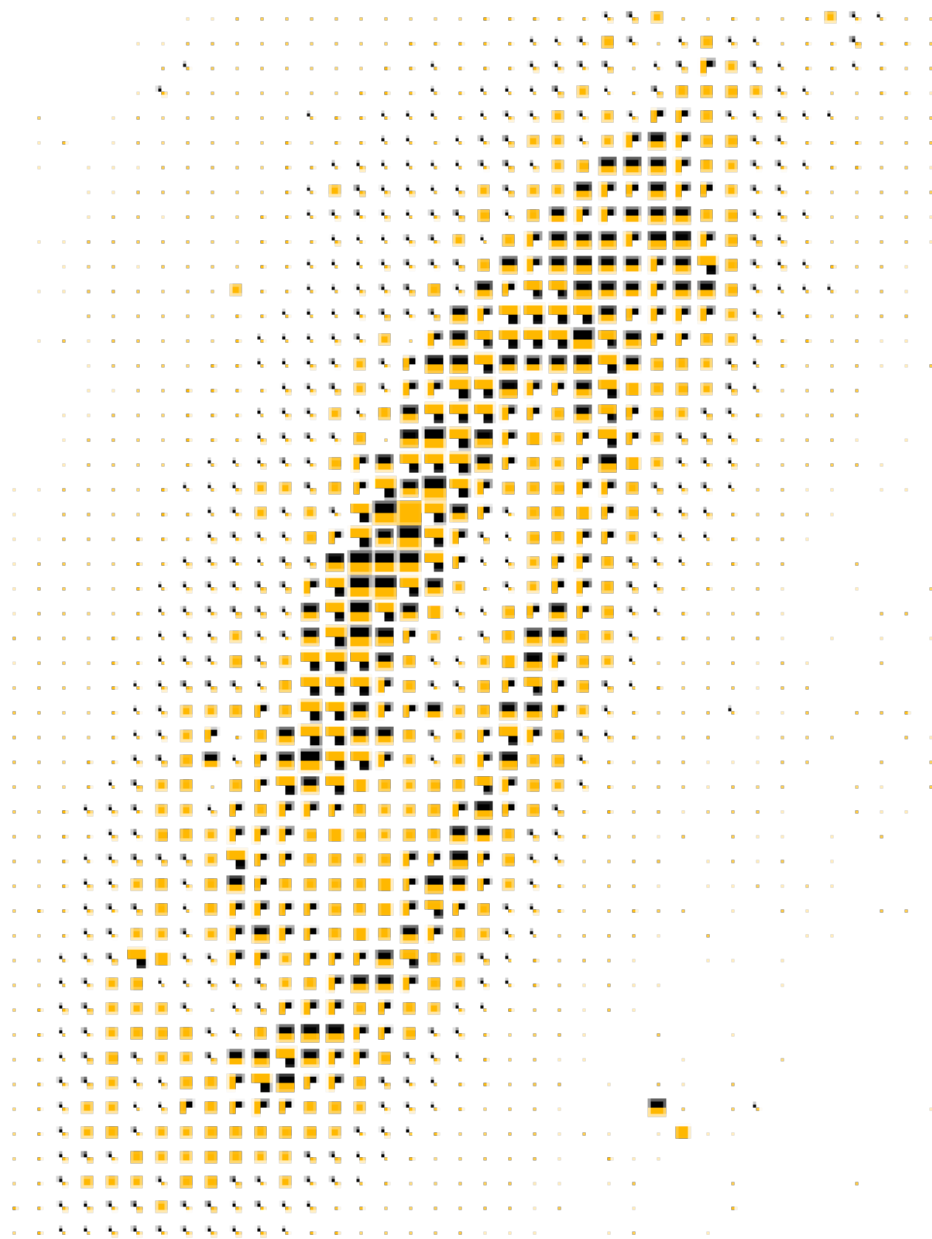


DND



# Delta Lake Performance Behind the Scenes: From Partitioning to Liquid Clustering

softserve



**DND**



## About me

- Big Data Engineer with extensive experience in Python
- Enthusiastic about math and machine learning
- Big fan of Remembrance of Earth's Past trilogy by Liu Cixin



**Roman Dryndik**

Senior Big Data Software Engineer

**softserve**



# **Agenda**

- 1. Intro**
- 2. Compaction**
- 3. Partitioning**
- 4. VACUUM**
- 5. Statistics and Data Skipping**
- 6. Z-Ordering**
- 7. Bloom Filter Index**
- 8. Liquid Clustering**
- 9. Summary**
- 10.Q&A**

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# Intro

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# Introduction

## Why Delta Lake Optimization Matters

### Key Problems:

- Small files problem: streaming writes, and DML operations generate thousands of tiny files → high metadata overhead, slow object-store listings, increased I/O costs, and poor compression
- Over-partitioning: high-cardinality partition keys (e.g., user\_id) create thousands of directories with micro-files → metadata explosion and degraded read performance
- Updates and deletes leave behind outdated files → table size grows and performance declines
- Inefficient data scanning: without proper data layout, queries scan terabytes instead of gigabytes



# Compaction

**Reducing Small Files to Speed Up Queries**



# Compaction

## The Small Files Problem

Root cause:

- **Streaming Ingestion:** writing micro-batches every few seconds/minutes
- **DML Operations:** frequent MERGE, UPDATE, or DELETE actions produce new files
- **Over-partitioning:** splitting data into too many granular folders



# Compaction

## The Small Files Problem

Performance impact:

- Growth of metadata in the transaction log
- **Metadata overhead:** the driver spends more time listing files in object storage (S3/ADLS) than processing data
- **High I/O latency:** opening/closing thousands of tiny files is inefficient
- **Poor compression:** parquet creates massive overhead (headers/footers) when files are too small





# Compaction

## OPTIMIZE & Auto Optimize

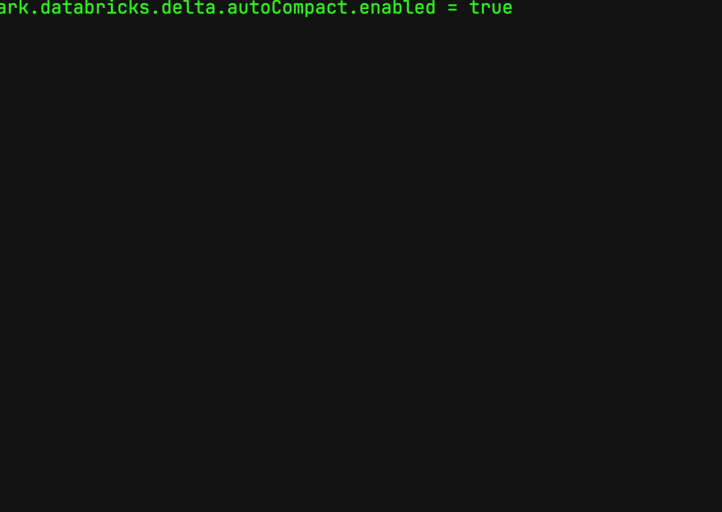
The **OPTIMIZE** command:

- Triggers **bin-packing**: reads small files and coalesces them into larger files (target size: **1 GB** by default)
- **Idempotent**: running it twice on the same data does nothing
- **ACID**: does not block concurrent readers or writers

**Auto Optimize** (Automated approach):

- **Optimized writes**: shuffles data **before** writing to reduce file count (increases write latency, improves read)
- **Auto compact**: triggers a "mini-optimize" **after** a write transaction commits

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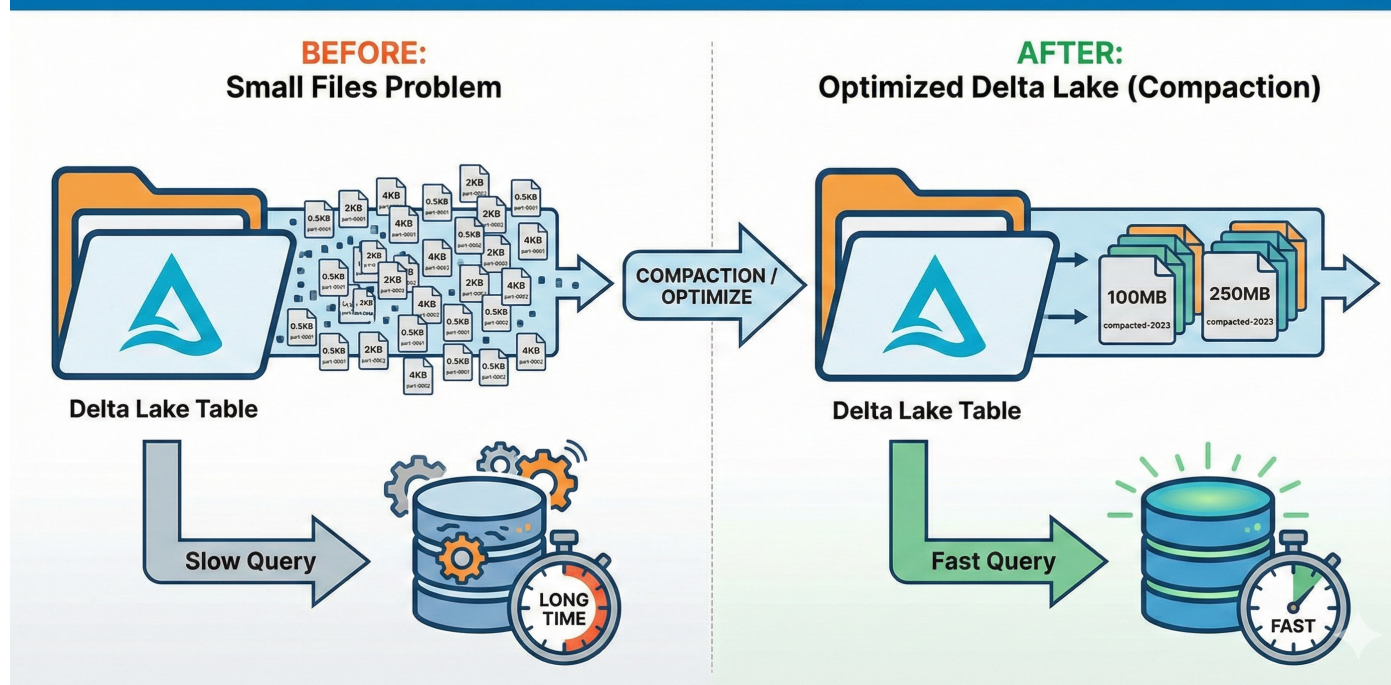


```
ss - vim compaction.py - vim - vim compaction.py - 80x24
1 spark.databricks.delta.optimizeWrite.enabled = true
2 spark.databricks.delta.autoCompact.enabled = true
~
~
~
~
~
~
~
~
~
~
~
~
~
~
~
~
~
~
```



# Compaction

## Delta Lake Small Files Problem





# Compaction

## Best Practices

- **Streaming:** enable **Auto Compact** + run OPTIMIZE periodically (e.g., daily/weekly) for cleanup
- **Batch ETL:** run OPTIMIZE at the end of the daily job
- **Strategy:** use **predicates** (WHERE) to avoid compacting the entire table every time



# Partitioning

**Splitting Data by Keys for Faster Queries**



# Partitioning

## Physical Data Layout

The concept:

- Splitting data into sub-directories based on high-level keys (e.g., Date, Region, Department)
- Structure: s3://bucket/table/date=2024-01-01/region=US/...

The benefit — partition pruning (aka data skipping):

- When a query includes a partition key in the WHERE clause, the engine completely ignores irrelevant directories
- Reduces scanned data from **terabytes** to **gigabytes**

The use case:

- Best for heavy filtering on specific columns (e.g., "Give me data for **Yesterday**").



# Partitioning

```
partitioning.sql

1  -- Create a table partitioned by date
2  CREATE TABLE sales (
3      id INT,
4      amount DOUBLE,
5      sale_date DATE
6  ) USING DELTA
7  PARTITIONED BY (sale_date);
8
9  -- Query that triggers Partition Pruning
10 SELECT sum(amount)
11 FROM sales
12 WHERE sale_date = '2024-01-01'; -- Skips all other
```



# Partitioning

## Over Partitioning

The trap — over-partitioning:

- Partitioning by high-cardinality columns (e.g., user\_id, timestamp, order\_id)
- **Result:** thousands of tiny directories containing tiny files
- **Impact:** severe metadata overhead (driver node bottleneck) and loss of compression efficiency

Typical symptoms:

- Many partitions containing **1–10 files**
- File sizes < 32 MB
- Query plan showing **hundreds of partitions scanned**
- OPTIMIZE is taking hours because the data is too fragmented





# Partitioning

## How to Avoid Over-Partitioning

General recommendations:

- Partition only by **low- or medium-cardinality** columns (e.g., date, country, category)
- Prefer **Z-ordering** or **Bloom filters** instead of “deep partitioning”
- If you have already over partitioned:
  - **Repartition** the table and rewrite it
  - Consolidate partitions (e.g., daily → monthly)

Rule of thumb:

- Partitions should generally contain at least **1 GB** of data
- If your table is small (< 1 TB), you might **not** need partitioning at all



# **VACUUM**

**Keeping Tables Clean and Performant**



# **VACUUM**

## **Problem**

- UPDATE, DELETE, INSERT, and MERGE operations create new files
- Old files increase the table size
- When your tables are large, you have poor read performance



# VACUUM

## Solution

- VACUUM operation — removes old unused parquet files



 vacuum.sql

```
1  -- Run VACUUM against table_name
2  VACUUM table_name [RETAIN num HOURS];
```



# VACUUM

## Solution

Retention period:

- Default: 7 days (RETAIN 168 HOURS)
- Can be reduced (e.g., for dev environments): RETAIN 24 HOURS
- Warning: too short → cannot time travel/rollback



 vacuum.sql

```
1  -- Run VACUUM against table_name
2  VACUUM table_name [RETAIN num HOURS];
```



# VACUUM

## Impact on Table Size & Performance

Effects of VACUUM:

- Reduces table size on disk
- Speeds up table scans
- Reduces cluster workload during reads

Tips:

- Run VACUUM after large updates/deletes
- Avoid running too frequently to preserve history



# **Statistics & Data Skipping**

## **Query Faster Without Full Scans**



# Statistics & Data Skipping

## File-Level Statistics

Delta Lake stores metadata for every data file:

- Min/max values per column
- Number of records
- Null counts
- Partition information (if applicable)

Why this matters:

- The client can understand the data distribution **without** opening the file
- Enables smarter, more targeted reading





# Statistics & Data Skipping

```
rdryn — admin@bozon: ~/projects/delta-lake-definitive-guide — ssh admin@192.168.0.10 — 123x32
{"fips":1001,"cases":0,"deaths":0,"maxValues":{"date":"2021-01-10","county":"Ziebach","state":"Wyoming"},
,"fips":78030,"cases":920560,"deaths":25562,"nullCount":{"date":0,"county":0,"state":0,"fips":1262,"cases":0,"deaths":3588}}
}
}
{
  "add": {
    "path": "part-00006-d0ec7722-b30c-4e1c-92cd-b4fe8d3bb954-c000.snappy.parquet",
    "partitionValues": {},
    "size": 883342,
    "modificationTime": 1619121488000,
    "dataChange": true,
    "stats": {"numRecords":147181,"minValues":{"date":"2021-01-10","county":"Abbeville","state":"Alabama",
,"fips":1001,"cases":0,"deaths":0,"maxValues":{"date":"2021-02-25","county":"Ziebach","state":"Wyoming"},
,"fips":78030,"cases":1188101,"deaths":29025,"nullCount":{"date":0,"county":0,"state":0,"fips":1250,"cases":0,"deaths":3510}}
}
}
{
  "add": {
    "path": "part-00007-4582392f-9fc2-41b0-ba97-a74b3afc8239-c000.snappy.parquet",
    "partitionValues": {},
    "size": 325440,
    "modificationTime": 1619121487000,
    "dataChange": true,
    "stats": {"numRecords":47559,"minValues":{"date":"2021-02-25","county":"Abbeville","state":"Alabama",
,"fips":1001,"cases":0,"deaths":0,"maxValues":{"date":"2021-03-11","county":"Ziebach","state":"Wyoming"},
,"fips":78030,"cases":1208672,"deaths":30068,"nullCount":{"date":0,"county":0,"state":0,"fips":408,"cases":0,"deaths":1170}}
}
}
}
admin@bozon:~/projects/delta-lake-definitive-guide$
```



# Statistics & Data Skipping

```
Zellij (fascinating-brachiosaur) Tab #1 Alt <[> VERTICAL
admin@bozon: ~/projects/delta-lake-definitive-guide — 0
  "dataChange": true,
  "stats": {"numRecords": 147181, "minValues": {"date":
: "2021-01-10", "county": "Abbeville", "state": "Alabama
", "fips": 1001, "cases": 0, "deaths": 0}, "maxValues": {"
date": "2021-02-25", "county": "Ziebach", "state": "Wyo
ming", "fips": 78030, "cases": 1188101, "deaths": 29025}, "
nullCount": {"date": 0, "county": 0, "state": 0, "fips": 12
50, "cases": 0, "deaths": 3510}}
}
}
{
  "add": {
    "path": "part-00007-4582392f-9fc2-41b0-ba97-a74b3afc8239
-c000.snappy.parquet",
    "partitionValues": {},
    "size": 325440,
    "modificationTime": 1619121487000,
    "dataChange": true,
    "stats": {"numRecords": 47559, "minValues": {"date":
\ "2021-02-25", "county": "Abbeville", "state": "Alabama\
", "fips": 1001, "cases": 0, "deaths": 0}, "maxValues": {"d
ate": "2021-03-11", "county": "Ziebach", "state": "Wyom
ing", "fips": 78030, "cases": 1208672, "deaths": 30068}, "n
ullCount": {"date": 0, "county": 0, "state": 0, "fips": 408
, "cases": 0, "deaths": 1170}}
}
}
admin@bozon:~/projects/delta-lake-definitive-guide$

admin@bozon: ~/projects/delta-lake-definitive-guide — 0
{
  "numRecords": 47559,
  "minValues": {
    "date": "2021-02-25",
    "county": "Abbeville",
    "state": "Alabama",
    "fips": 1001,
    "cases": 0,
    "deaths": 0
  },
  "maxValues": {
    "date": "2021-03-11",
    "county": "Ziebach",
    "state": "Wyoming",
    "fips": 78030,
    "cases": 1208672,
    "deaths": 30068
  },
  "nullCount": {
    "date": 0,
    "county": 0,
    "state": 0,
    "fips": 408,
    "cases": 0,
    "deaths": 1170
  }
}
admin@bozon:~/projects/delta-lake-definitive-guide$

Ctrl + <g> LOCK <p> PANE <t> TAB <n> RESIZE <h> MOVE <s> SEARCH <o> SESSION <q> QUIT ...
```



## Statistics & Data Skipping

**Data Skipping Means Client Avoids Reading Files That Cannot Satisfy the Query**

How it works:

- Client reads delta metadata (stats)
- Compares stats with query filters
- Only loads files whose min/max **overlaps** with the filter

Result:

- Fewer files scanned
- Less I/O
- Faster queries



## Statistics & Data Skipping

**Data Skipping Means Client Avoids Reading Files That Cannot Satisfy the Query**

What the client does:

- Checks file stats (min/max event\_date)
- Skips all files where **max** < '2023-06-01'
- Reads only relevant files instead of scanning the whole table

Impact:

- Significant reduction in scanned data
- Lower latency
- Better cluster utilization



filter.sql

```
1 SELECT *
2 FROM events
3 WHERE event_date ≥ '2023-06-01';
```



# **Z-Ordering**

## **Improving Multi-Column Filters**



# Z-Ordering

## What Is Z-Ordering?

Z-Ordering is a data clustering algorithm:

- Reorders data within Delta files
- Groups related column values close together on disk
- Uses a space-filling “Z-curve” to interleave bits from multiple columns

Goal:

- Make multi-column filtering faster by improving data locality



# Z-Ordering

## Why Z-Ordering?

Benefits:

- Reduces the number of files a client needs to scan
- More efficient data skipping
- Works best for **high-cardinality columns** (e.g., user\_id, timestamp)
- Particularly effective for multi-column predicates:



 z-curve.sql

```
1 WHERE user_id = X AND event_date BETWEEN ...
```



# Z-Ordering

## How Z-Order Helps

Without Z-order:

- Data for user\_id and event\_date is spread across many files, which means poor locality
- Client scans many files just to find a small subset of rows

With Z-order:

- Related values of user\_id + event\_date are co-located in fewer files
- The client can skip most files

Result:

- Significant improvements in selective queries





# Z-Ordering

## How Z-Order Helps

Scenario:

- You have a user events table

Columns:

- user\_id — high-cardinality identifier
- event\_date — timestamp/date
- event\_type — string (e.g., "click", "view", etc.)
- session\_id
- country
- device\_type

```
example.sql

1  SELECT *
2  FROM events
3  WHERE
4      user_id = '65e1419e-59db-4c5e-9914-b69361bac2fb' AND
5      event_date BETWEEN '2025-05-01' AND '2025-06-01';
```

Query pattern:

- Most queries filter by a combination of user\_id and event\_date



# Z-Ordering

## How Z-Order Helps

Without Z-order:

- Rows for the same user\_id are spread across many files
- Date ranges overlap across files unpredictably
- The client must scan **hundreds** of files

With Z-order on (user\_id, event\_date):

- Rows for each user\_id become physically close together
- Their dates fall into tight ranges
- The client can skip most files

```
z-curve-table-optimization-2.sql

1  -- Optimize table_name by user_id and event_date
2  OPTIMIZE events ZORDER BY (user_id, event_date);
```

Notes:

- Often combined with **OPTIMIZE** to compact small files and then apply Z-order
- Works best when you Z-order the columns most used together in filters
- Not necessary for partitioned columns



# Bloom Filter Index

**Fast Lookups on High-Cardinality Columns**



## Bloom Filter Index

### Why Bloom Filters?

Use case:

- Columns with **high cardinality** (UUIDs, emails, session IDs, order IDs)
- Traditional data skipping (min/max stats) doesn't help for these columns
- Bloom filters give a fast, space-efficient way to check “**might** contain this value?”

Result:

- Avoids scanning files that do **not** contain the value



# Bloom Filter Index

## Bloom Filter = Probabilistic Index

Benefits:

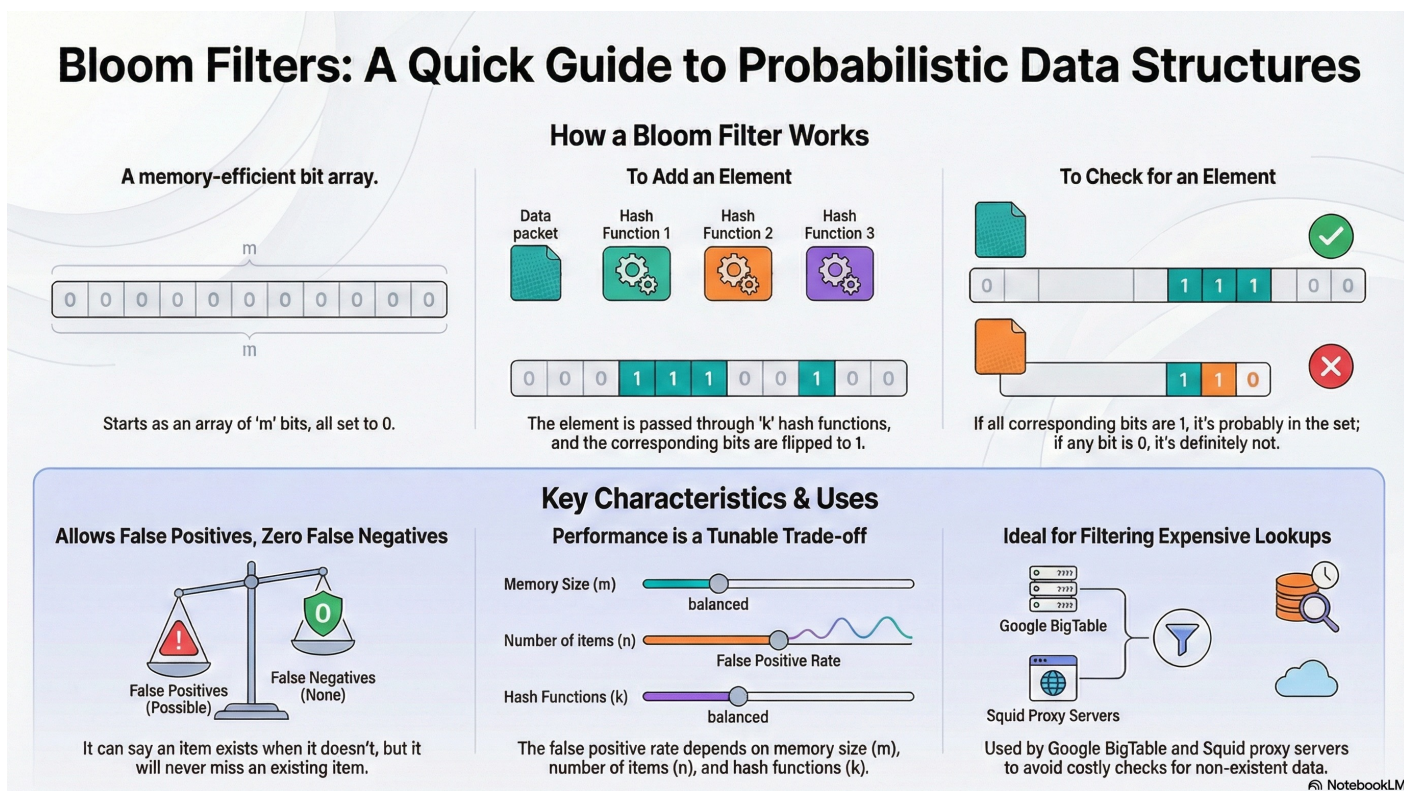
- Very compact bit array + multiple hash functions
- Can answer:
  - **“Definitely not present”**
  - **“Possibly present”** (false positives **possible**, false negatives **not** possible)

Why this helps:

- It is possible to check the Bloom filter **before** reading the file
- This eliminates large portion of irrelevant files → less I/O, faster reads



# Bloom Filter Index





## **Bloom Filter Index**

### **Example — UUID or Email Lookups**

Consider a table:

- user\_events

Columns:

- event\_id (UUID)
- user\_email (string)
- event\_timestamp
- event\_type



## Bloom Filter Index

### Example — UUID or Email Lookups

Without Bloom filters:

- UUIDs have no meaningful ordering
- min/max stats are useless
- Client may scan hundreds of files

```
bloom-filter.sql

1  SELECT *
2  FROM user_events
3  WHERE event_id = '550e8400-e29b-41d4-a716-446655440000';
```





## Bloom Filter Index

### Example — UUID or Email Lookups

With Bloom filter on event\_id:

```
bloom-filtrer-on-event-id.sql

1 ALTER TABLE user_events
2 ADD BLOOMFILTER INDEX
3 ON event_id
4 OPTIONS (fpp = 0.01, numItems = 1000000);
```



## Bloom Filter Index

### Example — UUID or Email Lookups

Effect:

- Client quickly checks each file's Bloom filter
- Most files are rejected immediately
- Only a few files are scanned

Moreover:

- Great for equality lookups on large, unique columns
- False positives are fine because it's still much faster than scanning
- Complements Z-order and statistics-based skipping
- **Small** index size means **low** overhead



# Liquid Clustering

## Adaptive Layout for Fast Queries



# Liquid Clustering

## Traditional Optimization Challenges

- Table management requires careful partition design
- Changing partition keys is expensive and often requires a full table rewrite
- Z-ordering needs continuous monitoring to match evolving query patterns
- Data skew: uneven distribution creates imbalanced partition sizes, which slows down queries



# Liquid Clustering

## What Is Liquid Clustering?

A new dynamic clustering strategy:

- Continuously maintains clustering as data evolves
- No need for expensive, full-table OPTIMIZE operations
- Automatically adapts as new data arrives or existing data changes
- It replaces partitioning and Z-ordering
- No more expensive rewrites

Idea:

- Instead of periodically “reorganizing everything,” the table stays well-organized over time
- The data is organized using clustering keys to optimize layout and simplify table management



# Liquid Clustering

## When to Use Liquid Clustering

- Frequent filtering on high-cardinality columns where partitioning fails
- Tables with data skew that need balanced distribution
- Fast-growing tables requiring constant tuning
- High-concurrency writes, where clustering reduces conflicts
- Changing query access patterns over time
- Cases where partitioning would create too many or too few partitions



# Liquid Clustering

## Enabling Liquid Clustering

- Add CLUSTER BY (<columns>) in CREATE TABLE for existing tables
- ALTER TABLE <name> CLUSTER BY (<columns>) for new tables
- Updates metadata only — does **not** rewrite existing data
- SQL, Python, and Scala APIs are all supported
- DataFrame API note: clustering keys can be set only at creation or with overwrite mode — not in append mode



# Liquid Clustering

## Enabling Liquid Clustering

Clustering is **incompatible** with traditional partitioning and ZORDER.  
It is designed to **replace** both.





# Liquid Clustering

## Choosing Clustering Keys

Selecting the right keys is crucial — good keys maximize data skipping and boost query performance.

- **Filter-first:** choose columns most used in WHERE clauses and joins
- **Stats required:** keys must be among columns with collected statistics (first 32 by default)
- **Avoid redundancy:** skip highly correlated columns — pick just one
- **Up to 4 keys:** more can hurt performance on tables < 10 TB



# Liquid Clustering

## Automatic Liquid Clustering

How it works:

- **Workload analysis:** identifies the most frequently filtered columns
- **Adaptive optimization:** updates keys as query patterns or data distribution change
- **Cost-aware decisions:** picks new keys only when the benefit outweighs the re-clustering cost
- **Powered by predictive optimization:** runs asynchronously in the background



clustering.sql

```
CREATE OR REPLACE TABLE table1(column01 int, column02 string)  
CLUSTER BY AUTO;
```



# Liquid Clustering

## Automatic Liquid Clustering

When keys may **not** be selected:

- The table is too small to benefit
- The existing layout is already effective
- Insufficient or inconsistent workload



# Liquid Clustering

## Key Benefits

- Stable performance without heavy maintenance jobs
- Faster queries with consistent data skipping
- Lower cost: avoids massive table rewrites
- Supports evolving datasets and dynamic clustering keys
- Redefine clustering keys without rewriting existing data is possible
- Simplify data layout: replaces manual partitioning and Z-order with one flexible technique
- Adapt to change: modify clustering keys without costly rewrites
- Automate with intelligence: automatic Liquid Clustering analyzes workloads and manages keys



## Unlock Faster Queries with Databricks Liquid Clustering

### The Modern Way to Organize Data



#### A Smarter Alternative to Partitioning & Z-Order

Liquid clustering automatically optimizes your data layout to speed up queries.

#### Evolve Your Data Layout On-the-Fly

You can redefine clustering keys at any time without rewriting existing data.



#### Ideal for Dynamic & Complex Tables

Best for tables with high-cardinality filters, skewed data, or changing access patterns.

### Getting Started with Liquid Clustering



#### Enable with a Simple `CLUSTER BY` Clause

Add ``CLUSTER BY (column_name)`` during table creation or when altering an existing table.



#### Let Databricks Do the Work with `CLUSTER BY AUTO`

This feature intelligently selects and adapts clustering keys based on your actual query workload.



#### Run `OPTIMIZE` to Apply Clustering

Periodically run the ``OPTIMIZE`` command to incrementally cluster new or updated data.

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# Summary

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## Summary & Recommendations

### Combine Techniques for Best Performance

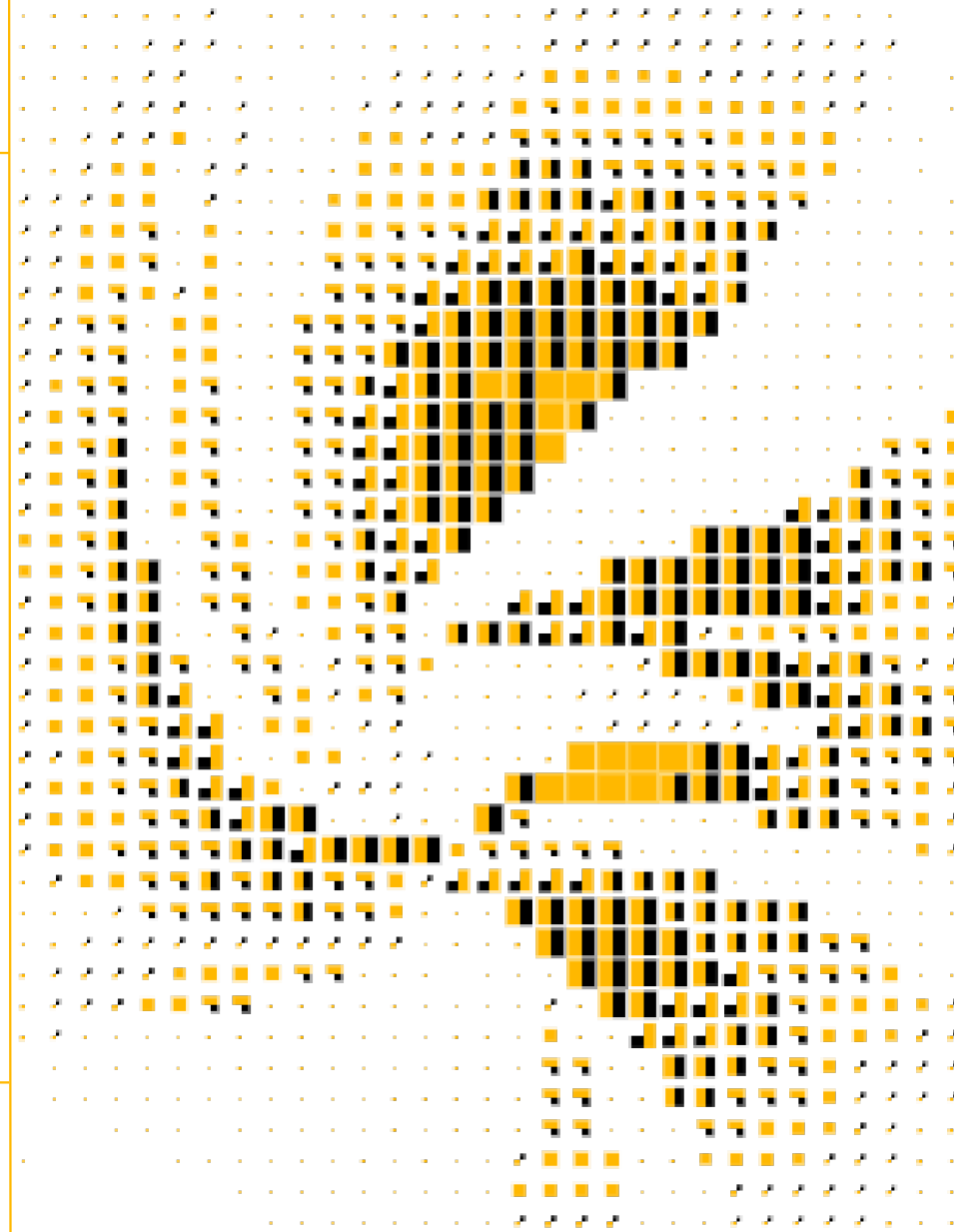
- No single optimization solves everything — Delta works best when techniques are combined
- Start simple: partition → compact → Z-order
- Add Bloom for faster lookups on high-cardinality columns
- Run **VACUUM** regularly
- Migrate to Liquid Clustering (if possible) for large, frequently-updated tables
- Re-evaluate periodically as data volume grows

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# Any Questions?

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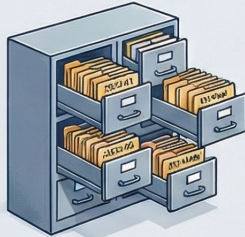


## Next-Gen Data Layout: Upgrading from Partitioning & Z-Ordering to Liquid Clustering

### The Old Way: Manual & Rigid Optimization

#### Partitioning: Coarse-Grained Pruning

Splits data into physical sub-directories based on low-cardinality keys (e.g., date, region).



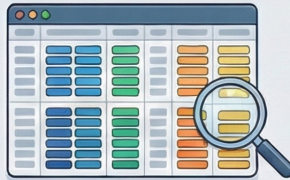
#### The "Over-Partitioning" Trap

Using high-cardinality keys (e.g., user, id) creates thousands of tiny files, crippling performance.



#### Z-Ordering: Static Data Co-location

Reorders data within files to group related column values together, improving multi-column filtering.



#### High Maintenance & Performance Decay

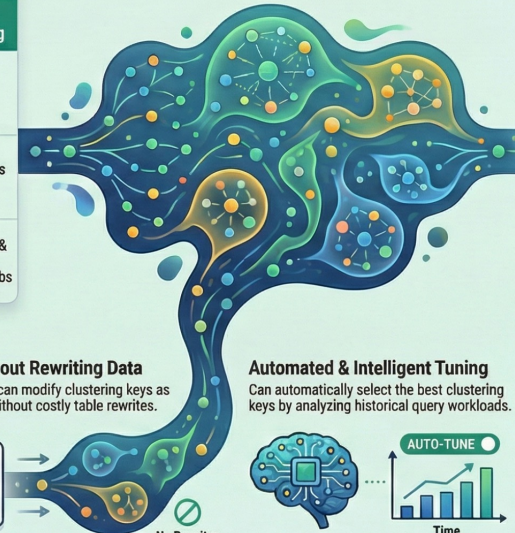
Requires periodic, expensive OPTIMIZE jobs and its effectiveness degrades as new data arrives.



### The New Way: Dynamic & Automated Optimization

#### Liquid Clustering: A Flexible, Adaptive Data Layout

Replaces both partitioning and Z-Ordering with a single, dynamic data organization strategy.

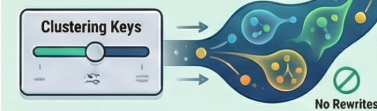


#### Feature Comparison

Feature	Partitioning	Z-Ordering	Liquid Clustering
Best For	Low-cardinality columns (date, country)	High-cardinality, multi-column filters	Evolving access patterns & data skew
Flexibility	Rigid (Full table rewrite to change keys)	Rigid (Static snapshot, needs re-running)	Flexible (Change keys without rewrites)
Maintenance	Manual tuning to avoid over-partitioning	Manual, periodic OPTIMIZE ZORDER jobs	Incremental & automated OPTIMIZE jobs

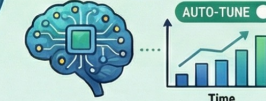
#### Change Keys Without Rewriting Data

Unlike partitioning, you can modify clustering keys as query patterns evolve without costly table rewrites.



#### Automated & Intelligent Tuning

Can automatically select the best clustering keys by analyzing historical query workloads.



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# **FOR THE FUTURE**

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